



LARN

Latency- and Resilience-Aware Networking

Latency- and Resilience-aware Transport for the Industrial Internet

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Internet (i.e. TCP/IP) is...

best-effort := no time-awareness but *fault-tolerant*.

Industrial Internet requires...

- ▶ ... fault-tolerance to cope with local failures.
- ▶ ... time-awareness to dependably interact with physical processes.

Industrial Internet protocols **MUST** provide time-awareness and reliability.

DFG SPP1914 “Cyber-Physical Networking”

- ▶ Joint research on **systems, network, and control engineering**.

“Latency- And Resilience-aware Networking” (Sep. 2016 - now)

- ▶ Telecommunications Chair at **Saarland Informatics Campus**
- ▶ Department of Computer Science 4 (Distributed Systems and Operating Systems) **Friedrich-Alexander-Universität Erlangen-Nürnberg**

State-of-the-Art (ε : Error, T : Latency)

- ▶ **(Simplified) Shannon:** $(\varepsilon \rightarrow 0) \Rightarrow (T = \infty)$
- ▶ **UDP** can neither make statements about ε nor bound T .
- ▶ **TCP** (and others with $\varepsilon = 0$) cannot bound T .

Industrial Internet & Multimedia Applications require bounds

- ▶ **Latency:** $\exists T_{max} : \forall T_i : T_i < T_{max} \wedge T_{max} \neq \infty$
- ▶ **Resilience:** $\varepsilon \approx 0$, but $\varepsilon \neq 0$

Goal: Develop an adaptive transport protocol that considers application requirements (latency and resilience bounds) and fulfills them if possible.

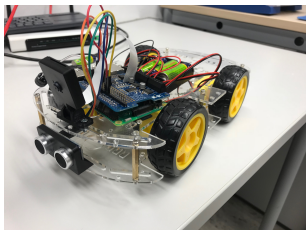
- ▶ **Reliable process-to-process real-time communication.**
 - Transport layer protocol.
- ▶ **Minimal requirements to lower layer.**
 - Rely on lower layer for forwarding, routing, process-multiplexing.
 - Allows to run on Ethernet / industrial busses.
- ▶ **No network knowledge required.**
 - Agnostic E2E mode.
 - Broad applicability, even on networks / links not under our control.
- ▶ **No network guarantees on channel parameters required.**
 - They are helpful, nevertheless.
- ▶ **Application must handle if constraints cannot be fulfilled currently.**
 - Channel is dynamic, so fulfilling constraints might be infeasible.

Measurement

- ▶ **Latency:** RTP-like
- ▶ **Loss:** Windowed packet tracking, Gilbert-Elliot model fitting
- ▶ **Data Rate:** Delivery Rate Estimation (IETF-Draft Cheng 2017)

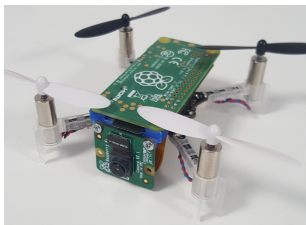
Control

- ▶ **Error Control:** Hybrid ARQ / Hybrid Error Coding.
- ▶ **Congestion Control:** BBR-inspired.
- ▶ **Rate Control:** Packet pacing.



Code

- ▶ Free Open Source Software
<http://prrt.larn.systems>
- ▶ C-Code + API (library)
- ▶ Python Bindings
- ▶ Gstreamer-Plugin



Platform Support

- ▶ **Requirement: glibc, pthreads**
- ▶ Raspberry Pi (3B, Zero W): Raspbian
- ▶ Odroid XU4: Ubuntu (16.04, 18.04)

Timing and **predictability** are crucial for the **Industrial Internet**.

Background

- ▶ Validation requires **per-packet timing** information.
- ▶ **Run-time evaluation necessary**, because ...
 - ▶ ... **hardware details** are not known in advance.
 - ▶ ... this information could be used to **adapt protocol operation at runtime**.

Idea: Instrument code with timestamping calls that have minimal impact on the results.

Solution: X-Lap

Cross-Layer Analysis Tool

(presented at Int. Workshop on Real-Time Networks (RTN17 & 18) at ECRTS)

Capturing

- ▶ Timestamping functions
- ▶ ⇒ Time- & cyclestamps
- ▶ C code

Analysis

- ▶ Data analysis
- ▶ ⇒ Latency and jitter
- ▶ Python code

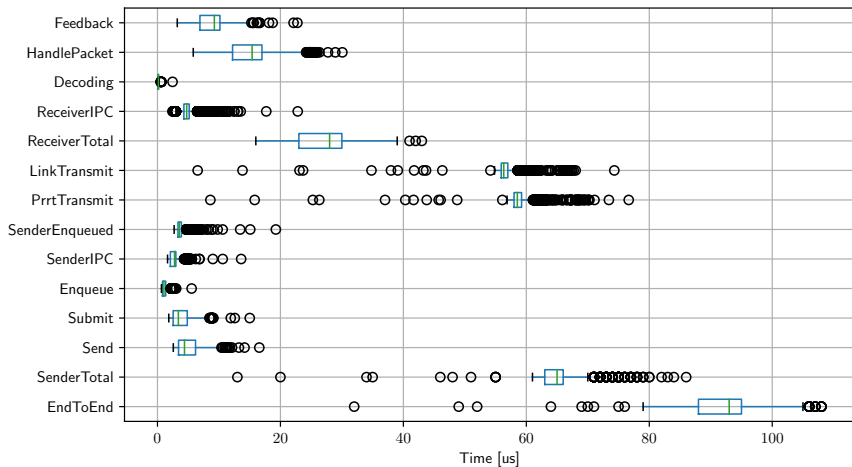
*.CSV



- ▶ Trace every packet
- ▶ Minimize run-time interference
- ▶ Embedded into protocol source code

- ▶ Single-packet traces
- ▶ Jitter amongst packets
- ▶ Latency criticality
- ▶ Correlation analysis

X-Lap: Trace Jitter



PRRT

- ▶ Predictably reliable real-time transport protocol.
- ▶ Error, congestion and rate control.
- ▶ Available for C, Python & Gstreamer.

For more information:
→ <http://larn.systems>

X-Lap

- ▶ Validation tool.
- ▶ Packet-level tracing & post-processing.
- ▶ Timing & resilience analysis.

Thank you for your attention. Questions?

Backup

General Approach (adapted re-implementation of Google's BBR)

- ▶ Measure...
 - ▶ T_{ch} (channel propagation delay) and
 - ▶ R_{ch} (channel bottleneck data rate).
- ▶ Keep $cwnd$ at one BDP ($T_{ch} \cdot R_{ch}$).
- ▶ Keep $pace$ between packets of size L by waiting for $\frac{L}{R_{ch}}$.

Adaptiveness

- ▶ Control sending $pace$ and $cwnd$ to do alternating measures of both channel parameters.
- ▶ Handle spurious loss events (caused by physical loss) differently to burst.

Approach not yet “perfect”, but controlling $cwnd$ and $pace$ is optimal strategy.

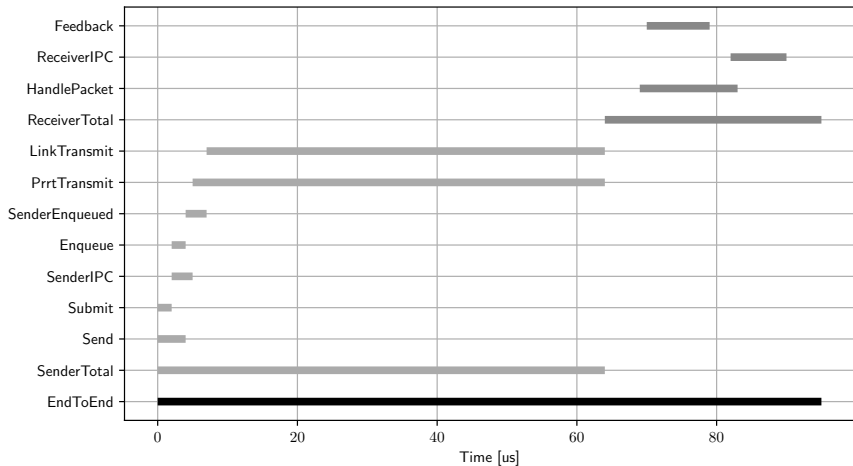
Hybrid Error Coding

- ▶ Configuration: $C = (n = 7, k = 4, \vec{n} = [1, 2])$

Adaptive Hybrid Error Coding

- ▶ Consider application constraints: $\varepsilon_{max}, T_{max}$.
- ▶ Measure channel parameters: $\varepsilon_{ch}, T_{ch}, R_{ch}$.
- ▶ Search for $C_{optimal}$ which...
 - ▶ ... fulfills application bounds.
 - ▶ ... minimizes caused redundancy (simplified: $\frac{n-k}{k}$).

X-Lap: Packet Traces



X-Lap: Correlation

